

Blocking Coloured Point Sets*

Greg Aloupis[†]Brad Ballinger[‡]Sébastien Collette[§]
David R. Wood^{**}Stefan Langerman[¶]Attila Pór^{||}

1 Introduction

This paper studies problems related to visibility and blocking in sets of coloured points in the plane. A point x *blocks* two points v and w if x is in the interior of the line segment \overline{vw} . Let P be a finite set of points in the plane. Two points v and w are *visible* with respect to P if no point in P blocks v and w . The *visibility graph* of P has vertex set P , where two distinct points $v, w \in P$ are adjacent if and only if they are visible with respect to P . A point set B *blocks* P if $P \cap B = \emptyset$ and for all distinct $v, w \in P$ there is a point in B that blocks v and w . That is, no two points in P are visible with respect to $P \cup B$, or alternatively, P is an independent set in the visibility graph of $P \cup B$.

A set of points P is *k-blocked* if each point in P is assigned one of k colours, such that each pair of points $v, w \in P$ are visible with respect to P if and only if v and w are coloured differently. Thus v and w are assigned the same colour if and only if some other point in P blocks v and w . We say P is $\{n_1, \dots, n_k\}$ -blocked if it is k -blocked and for some labelling of the colours by the integers $[k] := \{1, 2, \dots, k\}$, the i -th colour class has exactly n_i points, for each $i \in [k]$. Equivalently, P is $\{n_1, \dots, n_k\}$ -blocked if the visibility graph of P is the complete k -partite graph $K(n_1, \dots, n_k)$. See Figure 1 for an example.

The following fundamental conjecture regarding k -blocked point sets is the focus of this paper.

Conjecture 1 *For each integer k there is an integer n such that every k -blocked set has at most n points.*

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[†]Institute of Information Science, Academia Sinica, Taipei, Taiwan (aloupis.greg@gmail.com).

[‡]Department of Mathematics, Humboldt State University, Arcata, California, U.S.A (bjb86@humboldt.edu).

[§]Chargé de Recherches du F.R.S.-FNRS. Département d'Informatique, Université Libre de Bruxelles, Brussels, Belgium (sebastien.collette@ulb.ac.be).

[¶]Maitre de Recherches du F.R.S.-FNRS, Département d'Informatique, Université Libre de Bruxelles, Brussels, Belgium (stefan.langerman@ulb.ac.be).

^{||}Department of Mathematics, Western Kentucky University, Bowling Green, Kentucky, U.S.A. (attila.por@wku.edu).

**Dept. of Mathematics & Statistics, The University of Melbourne, Melbourne, Australia (woodd@unimelb.edu.au). QEII Research Fellow supported by the Australian Research Council.

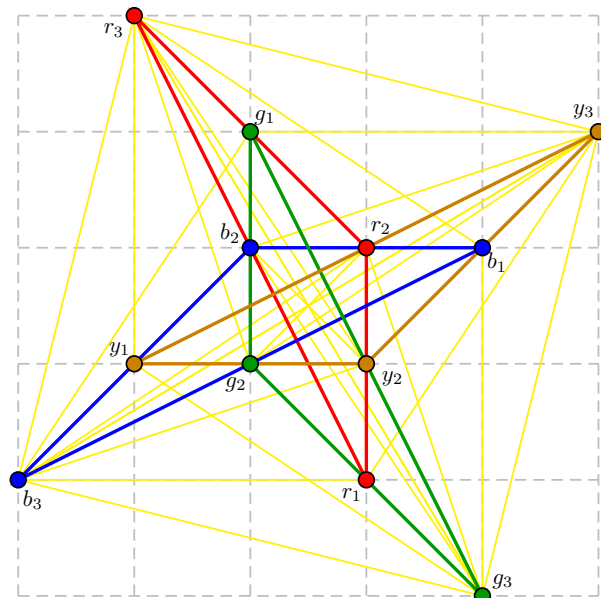


Figure 1: A $\{3, 3, 3\}$ -blocked point set.

A k -set is a multiset of k positive integers. A k -set $\{n_1, \dots, n_k\}$ is *representable* if there is an $\{n_1, \dots, n_k\}$ -blocked point set. As illustrated in Figure 2, it follows from the characterisation of 2- and 3-colourable visibility graphs by Kára et al. [6] that $\{1, 1\}$ and $\{1, 2\}$ are the only representable 2-sets, and that $\{1, 1, 1\}$, $\{1, 1, 2\}$, $\{1, 2, 2\}$ and $\{2, 2, 2\}$ are the only representable 3-sets. In particular, every 2-blocked point set has at most 3 points, and every 3-blocked point set has at most 6 points. This proves Conjecture 1 for $k \leq 3$. Later we prove Conjecture 1 for $k = 4$.

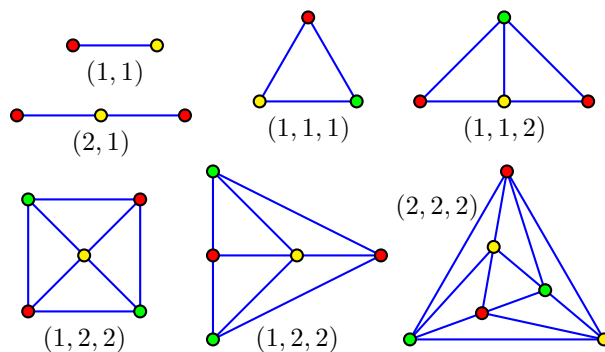


Figure 2: The 2-blocked and 3-blocked point sets.

This paper makes the following contributions. Section 2 introduces some background motivation. Section 3 describes methods for constructing k -blocked sets from a given $(k - 1)$ -blocked set. These methods lead to a characterisation of representable k -sets when each colour class has at most three points. Section 4 studies the $k = 4$ case in more detail. In particular, we characterise the representable 4-sets, and conclude that the example in Figure 1 is in fact the largest 4-blocked point set. Section 5 introduces a special class of k -blocked sets (so-called midpoint-blocked sets) that lead to a construction of the largest known k -blocked sets for infinitely many values of k .

Also note the following easily proved properties.

Lemma 1 ([2]) *At most three points are collinear in every k -blocked point set.*

Lemma 2 ([2]) *Each colour class in a k -blocked point set is in general position (no three collinear).*

2 Some Background Motivation

Much recent research on blockers began with the following conjecture by Kára et al. [6].

Conjecture 2 (Big-Line-Big-Clique Conjecture [6]) *For all integers t and ℓ there is an integer n such that for every finite set P of at least n points in the plane:*

- P contains ℓ collinear points, or
- P contains t pairwise visible points (that is, the visibility graph of P contains a t -clique).

Conjecture 2 is true for $t \leq 5$, but is open for $t \geq 6$ or $\ell \geq 4$; see [10, 1]. Jan Kára suggested the following weakening of Conjecture 2.

Conjecture 3 ([10]) *For all integers t and ℓ there is an integer n such that for every finite set P of at least n points in the plane:*

- P contains ℓ collinear points, or
- the chromatic number of the visibility graph of P is at least t .

Clearly Conjecture 2 implies Conjecture 3.

Proposition 3 *Conjecture 3 with $\ell = 4$ and $t = k + 1$ implies Conjecture 1.*

Proof. Assume Conjecture 3 holds for $\ell = 4$ and $t = k + 1$. Suppose P is a k -blocked set of at least n points. By Lemma 1, at most three points are collinear. Thus the first conclusion of Conjecture 3 does not hold. Since the visibility graph of P is k -colourable, the second conclusion of Conjecture 3 does not hold. This contradiction proves that every k -blocked set has less than n points, and Conjecture 1 holds. \square

Since Conjecture 2 holds for $t \leq 5$, Conjecture 1 holds for $k \leq 4$. Let $b(n)$ be the minimum integer such that some set of n points in the plane in general position is blocked by some set of $b(n)$ points. Linear lower bounds on $b(n)$ are known [7, 3], but many authors have conjectured or stated as an open problem that $b(n)$ is super-linear.

Conjecture 4 ([7, 9, 3, 10]) $\frac{b(n)}{n} \rightarrow \infty$ as $n \rightarrow \infty$.

Pór and Wood [10] proved that Conjecture 4 implies Conjecture 3, and thus implies Conjecture 1. That Conjecture 1 is implied by a number of other well-known conjectures, yet remains challenging, adds to its interest.

3 k -Blocked Sets with Small Colour Classes

We now describe some methods for building blocked point sets from smaller blocked point sets.

Lemma 4 *Let G be a visibility graph. Let $i \in \{1, 2, 3\}$. Furthermore suppose that if $i \geq 2$ then $V(G) \neq \emptyset$, and if $i = 3$ then not all the vertices of G are collinear. Let G_i be the graph obtained from G by adding an independent set of i new vertices, each adjacent to every vertex in G . Then G_1 , G_2 , and G_3 are visibility graphs.*

Proof. For distinct points p and q , let \overleftarrow{pq} denote the ray that is (1) contained in the line through p and q , (2) starting at p , and (3) not containing q . Let \mathcal{L} be the union of the set of lines containing at least two vertices in G .

$i = 1$: Since \mathcal{L} is the union of finitely many lines, there is a point $p \notin \mathcal{L}$. Thus p is visible from every vertex of G . By adding a new vertex at p , we obtain a representation of G_1 as a visibility graph.

$i = 2$: Let p be a point not in \mathcal{L} . Let v be a vertex of G . Each line in \mathcal{L} intersects \overleftarrow{vp} in at most one point. Thus $\overleftarrow{vp} \setminus \mathcal{L} \neq \emptyset$. Let q be a point in $\overleftarrow{vp} \setminus \mathcal{L}$. Thus p and q are visible from every vertex of G , but p and q are blocked by v . By adding new vertices at p and q , we obtain a representation of G_2 as a visibility graph.

$i = 3$: Let u, v, w be non-collinear vertices in G . Let p be a point not in \mathcal{L} and not in the convex hull of $\{u, v, w\}$. Without loss of generality, $\overline{uv} \cap \overline{pw} \neq \emptyset$. There are infinitely many pairs of points $q \in \overleftarrow{up}$ and $r \in \overleftarrow{vp}$ such that w blocks q and r . Thus there are such q and r both not in \mathcal{L} . By construction, u blocks p and q , and v blocks p and r . By adding new vertices at p , q and r , we obtain a representation of G_3 as a visibility graph. \square

We now characterise the representable (≥ 4)-sets, assuming each colour class has at most three points.

Proposition 5 A k -set $\{n_1, \dots, n_k\}$ is representable whenever $k \geq 4$ and each $n_i \leq 3$, except for $\{1, 3, 3, 3\}$ which is not representable [2].

Proof. We say $\{n_1, \dots, n_k\}$ contains $\{n_1, \dots, n_{i-1}, n_{i+1}, \dots, n_k\}$ for each $i \in [k]$. We proceed by induction on k . If $\{n_1, \dots, n_k\}$ contains a representable $(k-1)$ -set, then $\{n_1, \dots, n_k\}$ is also representable by Lemma 4. (Since $k \geq 4$ the assumptions in Lemma 4 hold.) Now assume that every $(k-1)$ -set contained in $\{n_1, \dots, n_k\}$ is not representable. By induction, we may assume that $k \leq 5$. Moreover, if $k = 5$ then $\{n_1, \dots, n_5\}$ must contain $\{1, 3, 3, 3\}$ (since by induction all other 4-sets are representable). Similarly, if $k = 4$ then $\{n_1, \dots, n_4\}$ must contain $\{1, 1, 3\}$, $\{1, 2, 3\}$, $\{1, 3, 3\}$, $\{2, 2, 3\}$, $\{2, 3, 3\}$ or $\{3, 3, 3\}$ (since $\{1, 1, 1\}$, $\{1, 1, 2\}$, $\{1, 2, 2\}$ and $\{2, 2, 2\}$ are representable). The following table describes the construction in each case.

$\{1, 1, 1, x\}$	contains $\{1, 1, 1\}$
$\{1, 1, 2, x\}$	contains $\{1, 1, 2\}$
$\{1, 1, 3, 3\}$	Figure 1 minus $\{r_1, g_3, r_3, g_1\}$
$\{1, 2, 2, x\}$	contains $\{1, 2, 2\}$
$\{1, 2, 3, 3\}$	Figure 1 minus $\{g_1, g_3, r_3\}$
$\{2, 2, 2, x\}$	contains $\{2, 2, 2\}$
$\{2, 2, 3, 3\}$	Figure 1 minus $\{g_3, r_3\}$
$\{2, 3, 3, 3\}$	Figure 1 minus g_3
$\{1, 1, 3, 3, 3\}$	contains $\{1, 1, 3, 3\}$
$\{1, 2, 3, 3, 3\}$	contains $\{1, 2, 3, 3\}$
$\{1, 3, 3, 3, 3\}$	contains $\{3, 3, 3, 3\}$

□

4 4-Blocked Point Sets

As shown in Section 2, Conjecture 1 holds for $k \leq 4$. That is, every 4-blocked set has bounded size. An explicit bound of 2^{790} follows from a result of Abel et al. [1], which can be improved to 2^{578} using a recent result by Dumitrescu et al. [3]; see [2]. Before characterising all representable 4-sets, we give a simple proof that every 4-blocked point set has bounded size.

Proposition 6 Every 4-blocked set has at most 36 points.

Proof. Let P be a 4-blocked set. Suppose that $|P| \geq 37$. Let S be the largest colour class. Thus $|S| \geq 10$. By Lemma 2, S is in general position. By a theorem of Harborth [4], some 5-point subset $K \subseteq S$ is the vertex-set of an empty convex pentagon $\text{conv}(K)$. Let $T := P \cap (\text{conv}(K) - K)$. Since $\text{conv}(K)$ is empty with respect to S , each point in T is not in S . Thus T is 3-blocked. K needs at least 8 blockers (5 blockers for the edges on the boundary of $\text{conv}(K)$, and 3 blockers for the chords of $\text{conv}(K)$). Thus $|T| \geq 8$. But every 3-blocked set has at most 6 points, which is a contradiction. Hence $|P| \leq 36$. □

Theorem 7 A 4-set $\{a, b, c, d\}$ is representable if and only if:

- $\{a, b, c, d\} = (4, 2, 2, 1)$, or
- $\{a, b, c, d\} = (4, 2, 2, 2)$, or
- all of $a, b, c, d \leq 3$ except for $\{3, 3, 3, 1\}$

Proof Sketch. Figure 3 shows $\{4, 2, 2, 1\}$ -blocked and $\{4, 2, 2, 2\}$ -blocked point sets. When $a, b, c, d \leq 3$, the required constructions are described in Proposition 5. Now we prove the converse. Let P be a 4-blocked point set. We prove [2] that if some colour class S contains a 4-point subset K , such that $\text{conv}(K)$ is a convex quadrilateral that is empty with respect to S , then P is $\{4, 2, 2, 1\}$ -blocked. Moreover, if some colour class S has at least five points, then by Lemma 2 and a theorem of Esther Klein, S contains such a subset K —implying P is $\{4, 2, 2, 1\}$ -blocked, which is a contradiction. Hence each colour class has at most four points. Let S be a largest colour class. If S consists of four points in convex position, then P is $\{4, 2, 2, 1\}$ -blocked (just set $K := S$). If S consists of four points in nonconvex position, then we prove [2] that P is $\{4, 2, 2, 2\}$ -blocked. Otherwise $|S| \leq 3$, and we are done by Proposition 5. □

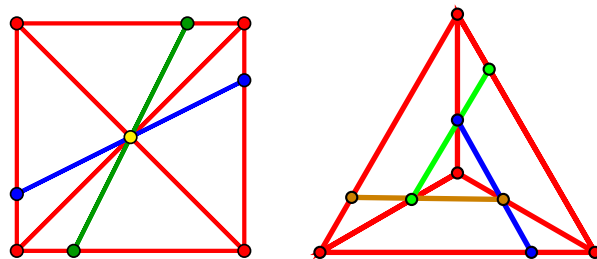


Figure 3: $\{4, 2, 2, 1\}$ -blocked and $\{4, 2, 2, 2\}$ -blocked point sets.

Corollary 8 Every 4-blocked set has at most 12 points, and there is a 4-blocked set with 12 points.

5 Midpoint-Blocked Point Sets

A k -blocked point set P is k -midpoint-blocked if for each monochromatic pair of distinct points $v, w \in P$ the midpoint of \overline{vw} is in P . Of course, the midpoint of \overline{vw} blocks v and w . A point set P is $\{n_1, \dots, n_k\}$ -midpoint-blocked if it is $\{n_1, \dots, n_k\}$ -blocked and k -midpoint-blocked. For example, the point set in Figure 1 is $\{3, 3, 3, 3\}$ -midpoint-blocked.

Another interesting example is the projection¹ of $[3]^d$. With $d = 1$ this point set is $\{2, 1\}$ -blocked, with $d = 2$ it is $\{4, 2, 2, 1\}$ -blocked, and with $d = 3$ it is $\{8, 4, 4, 4, 2, 2, 2, 1\}$ -blocked. In general, each set of

¹If G is the visibility graph of some point set $P \subseteq \mathbb{R}^d$, then G is the visibility graph of some projection of P to \mathbb{R}^2 .

points with exactly the same set of coordinates equal to 2 is a colour class, there are 2^{d-i} colour classes of points with exactly i coordinates equal to 2, and $[3]^d$ is $\{\binom{d}{i} \times 2^i : i \in [0, d]\}$ -midpoint-blocked and 2^d -midpoint-blocked.

Hernández-Barrera et al. [5] defined $m(n)$ to be the minimum number of midpoints determined by some set of n points in general position in the plane, and proved that $m(n) \leq cn^{\log_2 3} = cn^{1.585\dots}$. This upper bound was improved by Pach [8] (and later by Matousek [7]) to $m(n) \leq nc^{\sqrt{\log n}}$. Hernández-Barrera et al. [5] conjectured that $m(n)$ is super-linear, which was verified by Pach [8]; that is, $\frac{m(n)}{n} \rightarrow \infty$ as $n \rightarrow \infty$. Pór and Wood [10] proved the following more precise version: For some $c > 0$, for all $\epsilon > 0$ there is an integer $N(\epsilon)$ such that $m(n) \geq cn(\log n)^{1/(3+\epsilon)}$ for all $n \geq N(\epsilon)$.

Theorem 9 *For each k there is an integer n such that every k -midpoint-blocked set has at most n points. More precisely, there is an absolute constant c and for each $\epsilon > 0$ there is an integer $N(\epsilon)$, such that for all k , every k -midpoint-blocked set has at most $k \max\{N(\epsilon), c^{(k-1)^{3+\epsilon}}\}$ points.*

Proof. Let P be k -midpoint-blocked set of n points. We may assume that $\frac{n}{k} > N(\epsilon)$. Let S be a set of exactly $s := \lceil \frac{n}{k} \rceil$ monochromatic points in P . Thus S is in general position by Lemma 2. And for every pair of distinct points $v, w \in S$ the midpoint of \overline{vw} is in $P - S$. Thus $c \frac{n}{k} (\log \frac{n}{k})^{1/(3+\epsilon)} \leq m(s) \leq n - s \leq n(1 - \frac{1}{k})$. Hence $(\log \frac{n}{k})^{1/(3+\epsilon)} \leq (k-1)/c$, implying $n \leq k2^{((k-1)/c)^{3+\epsilon}}$. The result follows. \square

We now construct k -midpoint-blocked point sets with a ‘large’ number of points. The method is based on the following product of point sets P and Q . Let (x_v, y_v) be the coordinates of each $v \in P \cup Q$. Let $P \times Q$ be the point set $\{(v, w) : v \in P, w \in Q\}$ where (v, w) is at (x_v, y_v, x_w, y_w) in 4-dimensional space. For brevity we do not distinguish between a point in \mathbb{R}^4 and its image in an occlusion-free projection of the visibility graph of $P \times Q$ into \mathbb{R}^2 .

Lemma 10 *If P is a $\{n_1, \dots, n_k\}$ -midpoint-blocked point set and Q is a $\{m_1, \dots, m_\ell\}$ -midpoint-blocked point set, then $P \times Q$ is $\{n_i m_j : i \in [k], j \in [\ell]\}$ -midpoint-blocked.*

Proof. Colour each point (v, w) in $P \times Q$ by the pair $(\text{col}(v), \text{col}(w))$. There are $n_i m_j$ points for the (i, j) -th colour class. It is straightforward to verify that two points in $P \times Q$ are blocked if and only if they have the same colour. Thus $P \times Q$ is blocked. Since every blocker is a midpoint, $P \times Q$ is midpoint-blocked. \square

Say P is a k -midpoint blocked set of n points. By Lemma 10, the i -fold product $P^i := P \times \dots \times P$ is a k^i -blocked set of $n^i = (k^i)^{\log_k n}$ points. Taking P to be

the $\{3, 3, 3, 3\}$ -midpoint-blocked point set in Figure 1, we obtain the following result, which describes the largest known construction of k -blocked point sets.

Theorem 11 *For all k a power of 4, there is a k -blocked set of $k^{\log_4 12} = k^{1.79\dots}$ points.*

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