COMP 2804 — Solutions Assignment 2

Question 1: On the first page of your assignment, write your name and student number.

Solution:

• Name: Sidney Crosby

• Student number: 87

Question 2: The function $f: \mathbb{N} \to \mathbb{N}$ is defined by

$$f(0) = 1,$$

 $f(n) = \frac{1}{2} \cdot 4^n \cdot f(n-1)$ if $n \ge 1$.

Prove that for every integer $n \geq 0$,

$$f(n) = 2^{n^2};$$

this reads as 2 to the power n^2 .

Solution: The proof is by induction on n. The base case is when n=0. Since f(0)=1 and

$$2^{n^2} = 2^{0^2} = 2^0 = 1,$$

the base case holds.

Let $n \ge 1$ and assume that the claim is true for n-1. Thus, the induction hypothesis is that

$$f(n-1) = 2^{(n-1)^2}.$$

We have to show that

$$f(n) = 2^{n^2}.$$

Using the recurrence, the induction hypothesis, and some basic algebra, we get

$$f(n) = \frac{1}{2} \cdot 4^{n} \cdot f(n-1)$$

$$= 2^{-1} \cdot 2^{2n} \cdot 2^{(n-1)^{2}}$$

$$= 2^{-1} \cdot 2^{2n} \cdot 2^{n^{2}-2n+1}$$

$$= 2^{n^{2}}.$$

Question 3: The functions $f: \mathbb{N} \to \mathbb{N}$ and $g: \mathbb{N}^2 \to \mathbb{N}$ are recursively defined as follows:

$$f(0) = 1,$$

 $f(n) = g(f(n-1), 2n)$ if $n \ge 1,$
 $g(0,n) = 0$ if $n \ge 0,$
 $g(m,n) = g(m-1,n) + n$ if $m \ge 1$ and $n \ge 0.$

Solve these recurrence relations for f, i.e., express f(n) in terms of n. Justify your answer. Hint: Start by solving the recurrence relation for g.

Solution: If you stare long enough at the recurrence for g, it makes sense that

$$g(m,n) = mn$$

for all $m \ge 0$ and $n \ge 0$. We prove by induction on m that this is indeed the case.

The base case is when m = 0. Since g(0, n) = 0 and $mn = 0 \cdot n = 0$, the base case holds. Let $m \ge 1$ and assume that

$$g(m-1,n) = (m-1)n.$$

Then

$$g(m,n) = g(m-1,n) + n = (m-1)n + n = mn.$$

Now that we have solved the recurrence for g, we can rewrite the recurrence for f:

$$f(0) = 1,$$

 $f(n) = 2n \cdot f(n-1)$ if $n \ge 1$.

The recursive rule says: To get f(n), take the previous value f(n-1), and multiply it by 2 and by n. From this, it makes sense to guess that

$$f(n) = 2^n \cdot n!$$

for all $n \geq 0$. We prove by induction on n that this is indeed the case.

The base case is when n=0. Since f(0)=1 and

$$2^n \cdot n! = 2^0 \cdot 0! = 1 \cdot 1 = 1,$$

the base case holds.

Let $n \geq 1$ and assume that the claim is true for n-1. Thus, we assume that

$$f(n-1) = 2^{n-1} \cdot (n-1)!.$$

Then we get

$$f(n) = 2n \cdot f(n-1)$$

$$= 2n \cdot 2^{n-1} \cdot (n-1)!$$

$$= (2 \cdot 2^{n-1}) (n \cdot (n-1)!)$$

$$= 2^n \cdot n!.$$

Question 4: For any integer $n \ge 1$, a permutation a_1, a_2, \ldots, a_n of the set $\{1, 2, \ldots, n\}$ is called *awesome*, if the following condition holds:

• For every i with $1 \le i \le n$, the element a_i in the permutation belongs to the set $\{i-1, i, i+1\}$.

For example, for n = 5, the permutation 2, 1, 3, 5, 4 is awesome, whereas 2, 1, 5, 3, 4 is not an awesome permutation.

Let P_n denote the number of awesome permutations of the set $\{1, 2, \ldots, n\}$.

- Determine P_1 , P_2 , and P_3 .
- Determine the value of P_n , i.e., express P_n in terms of numbers that we have seen in class. Justify your answer.

Hint: Derive a recurrence relation. What are the possible values for the last element a_n in an awesome permutation?

Solution:

- n = 1: There is only one permutation of the set $\{1\}$, namely 1. This permutation is awesome and, therefore, $P_1 = 1$.
- n = 2: There are two permutations of the set $\{1, 2\}$, namely 12 and 21. Both are awesome and, therefore, $P_2 = 2$.
- n = 3: If a permutation $a_1a_2a_3$ of the set $\{1, 2, 3\}$ is awesome, then $a_1 \in \{1, 2\}$ and $a_3 \in \{2, 3\}$. This leads to three awesome permutations: 123, 132, and 213. Thus, $P_3 = 3$.

If you want, you can consider the other three permutations 231, 312, and 321 and convince yourself that neither of these is awesome.

- Let $n \ge 3$ and consider an awesome permutation a_1, a_2, \ldots, a_n of the set $\{1, 2, \ldots, n\}$. We follow the hint: The value a_n can be either n-1 or n.
 - Assume $a_n = n$. Then $a_1, a_2, \ldots, a_{n-1}$ is an awesome permutation of the set $\{1, 2, \ldots, n-1\}$. There are P_{n-1} many permutations of this type.
 - Assume $a_n = n 1$. Then a_{n-1} must be equal to n. (Otherwise, there is an i with $1 \le i \le n 2$ such that $a_i = n$. But then the permutation is not awesome.) Then $a_1, a_2, \ldots, a_{n-2}$ is an awesome permutation of the set $\{1, 2, \ldots, n-2\}$. There are P_{n-2} many permutations of this type.

Conclusion: On the one hand, the number of awesome permutations is equal to P_n . On the other hand, the number of such permutations is equal to $P_{n-1} + P_{n-2}$.

• We have obtained the recurrence $P_1 = 1$, $P_2 = 2$, and $P_n = P_{n-1} + P_{n-2}$ for $n \ge 3$. This is a shifted Fibonacci sequence and it follows that $P_n = f_{n+1}$ for all $n \ge 1$. Question 5: The Fibonacci numbers are defined as follows: $f_0 = 0$, $f_1 = 1$, and $f_n = f_{n-1} + f_{n-2}$ for $n \ge 2$.

In class, we have seen that for any $m \ge 1$, the number of 00-free bitstrings of length m is equal to f_{m+2} . (In class, I showed this for $m \ge 2$, but this result is also valid for m = 1.) Let $n \ge 1$ be an integer. For each question below, justify your answer.

• How many 00-free bitstrings of length n+2 do not contain any 0?

Solution: Such a string contains only 1's. There is only one such string. Thus, the answer is

$$1. (1)$$

• How many 00-free bitstrings of length n+2 contain exactly one 0?

Solution: Such a string contains one 0 and n + 1 many 1's. Since there are n + 2 positions for the bit 0, the answer is

$$n+2. (2)$$

• How many 00-free bitstrings of length n+2 have the following property: The bitstring contains at least two 0's, and the second rightmost 0 is at position 1.

Solution: Such a string must start with 01 and it contains exactly one 0 in the positions $3, 4, \ldots, n+2$. Since there are n positions for this bit 0, the answer is

$$n = n \cdot f_1. \tag{3}$$

• How many 00-free bitstrings of length n+2 have the following property: The bitstring contains at least two 0's, and the second rightmost 0 is at position 2.

Solution: Such a string must start with 101 and it contains exactly one 0 in the positions $4, 5, \ldots, n+2$. Since there are n-1 positions for this bit 0, the answer is

$$n - 1 = (n - 1) \cdot f_2. \tag{4}$$

• Let k be an integer with $3 \le k \le n$. How many 00-free bitstrings of length n+2 have the following property: The bitstring contains at least two 0's, and the second rightmost 0 is at position k.

Solution: Such a string can be divided into three pieces:

- The middle piece is the substring at positions k-1, k, k+1. We know that there is a 0 at position k. Since the entire string is 00-free, there is a 1 at position k-1, and a 1 at position k+1. Thus, this middle piece is 101.

- The left piece is the substring at positions 1, 2, ..., k-2. This left piece can be any 00-free bitstring of length k-2. Thus, there are f_k many possibilities for this left piece.
- The right piece is the substring at positions $k+2, k+3, \ldots, n+2$; this piece has length n-k+1. In this right piece, there is exactly one 0. This 0 can be in any of the n-k+1 possible positions. Thus, there are n-k+1 many possibilities for this right piece.
- By the Product Rule, the answer to this part of the question is

$$(n-k+1)\cdot f_k. (5)$$

• Let k be an element of $\{n+1, n+2\}$. How many 00-free bitstrings of length n+2 have the following property: The bitstring contains at least two 0's, and the second rightmost 0 is at position k.

Solution: If k = n+1, then the second rightmost 0 cannot be at position k; otherwise, the rightmost 0 is at position k+1=n+2, and the string is not 00-free.

If k = n + 2, then the second rightmost 0 cannot be at position k; otherwise, there is no space for the rightmost 0.

Thus, the answer to this part of the question is

$$0. (6)$$

• Use the previous results to prove that

$$\sum_{k=1}^{n} (n-k+1) \cdot f_k = f_{n+4} - n - 3,$$

i.e.,

$$n \cdot f_1 + (n-1) \cdot f_2 + (n-2) \cdot f_3 + \dots + 2 \cdot f_{n-1} + 1 \cdot f_n = f_{n+4} - n - 3.$$

Solution: We know that the number of 00-free bitstrings of length n + 2 is equal to f_{n+4} . Each such string is of exactly one of the types as we have considered above. Thus, the sum of (1), (2), (3), (4), (5), and (6) is equal to f_{n+4} .

Question 6: Those of you who come to class will remember that Elisa Kazan¹ loves to drink cider. After a week of bossing the Vice-Presidents around, Elisa goes to the pub and runs the following recursive algorithm, which takes as input an integer $n \ge 0$:

¹President of the Carleton Computer Science Society

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Algorithm ELISAGOESTOTHEPUB(n):

if n = 0
then drink one bottle of cider
else for k = 0 to n - 1
do ELISAGOESTOTHEPUB(k);
drink one bottle of cider
endfor
endif
```

For $n \geq 0$, let C(n) be the number of bottles of cider that Elisa drinks when running algorithm ELISAGOESTOTHEPUB(n).

Prove that for every integer $n \geq 1$,

$$C(n) = 3 \cdot 2^{n-1} - 1.$$

Hint:
$$1 + 2 + 2^2 + 2^3 + \dots + 2^{n-2} = 2^{n-1} - 1$$
.

Solution: If Elisa runs ELISAGOESTOTHEPUB(0), then she drinks one bottle of cider. Thus,

$$C(0) = 1.$$

Let $n \geq 1$ and consider what happens when Elisa runs ELISAGOESTOTHEPUB(n). The for-loop makes n iterations, one for every k = 0, 1, 2, ..., n - 1. In the k-th iteration, (i) Elisa runs ELISAGOESTOTHEPUB(k), during which she drinks C(k) bottles of cider, and (ii) Elisa drinks one bottle of cider. Overall, in the k-th iteration, Elisa drinks 1 + C(k) bottles of cider. We conclude that, for $n \geq 1$,

$$C(n) = \sum_{k=0}^{n-1} (1 + C(k)).$$

This is the same as

$$C(n) = n + \sum_{k=0}^{n-1} C(k).$$

In words, C(0) = 1. For any $n \ge 1$, to obtain C(n), we take the sum of n and the total sum of all previous C-values.

It remains to verify that this recurrence relation solves to

$$C(n) = 3 \cdot 2^{n-1} - 1,$$

for each $n \ge 1$. (Note that this is not true for n = 0.) We prove this by induction:

The base case is when n = 1. Since

$$C(1) = 1 + C(0) = 1 + 1 = 2$$

and

$$3 \cdot 2^{n-1} - 1 = 3 \cdot 2^0 - 1 = 3 \cdot 1 - 1 = 2$$

the base case holds.

Let $n \geq 2$ and assume that for all $1 \leq k \leq n-1$,

$$C(k) = 3 \cdot 2^{k-1} - 1.$$

Then, using the recurrence for C(n), the induction hypothesis, and the hint, we get

$$C(n) = n + \sum_{k=0}^{n-1} C(k)$$

$$= n + C(0) + \sum_{k=1}^{n-1} C(k)$$

$$= n + 1 + \sum_{k=1}^{n-1} (3 \cdot 2^{k-1} - 1)$$

$$= n + 1 + \left(3 \sum_{k=1}^{n-1} 2^{k-1}\right) - (n-1)$$

$$= 2 + \left(3 \sum_{k=1}^{n-1} 2^{k-1}\right)$$

$$= 2 + 3 \cdot (1 + 2 + 2^2 + 2^3 + \dots + 2^{n-2})$$

$$= 2 + 3 \cdot (2^{n-1} - 1)$$

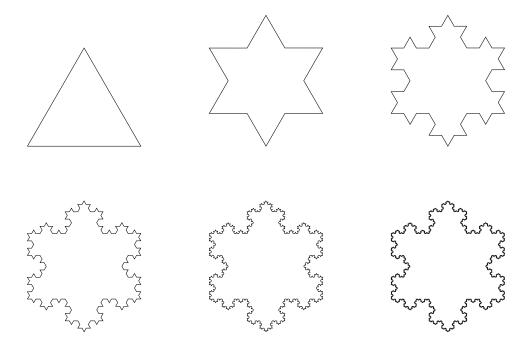
$$= 3 \cdot 2^{n-1} - 1.$$

Question 7: The sequence SF_0, SF_1, SF_2, \ldots of *snowflakes* is recursively defined in the following way:

- The snowflake SF_0 is an equilateral triangle with edges of length 1.
- For any integer $n \geq 1$, the snowflake SF_n is obtained by taking the snowflake SF_{n-1} and doing the following for each of its edges:
 - Divide this edge into three edges of equal length.
 - Draw an equilateral triangle that has the middle edge from the previous step as its base, and that is outside of SF_{n-1} .
 - Remove the edge that is the base of the equilateral triangle from the previous step.

Note: In the original question, these were called crystals. I changed the name to snowflakes, because that is what these things are called.

In the figure below, you see the snowflakes SF_0 up to SF_5 .



• For any integer $n \geq 0$, let N_n be the total number of edges of SF_n . Determine the value of N_n , by deriving a recurrence relation and solving it.

Solution: Since SF_0 is a triangle, we have $N_0 = 3$. Let $n \ge 1$. We obtain SF_n by replacing each edge in SF_{n-1} by four edges. This implies that $N_n = 4 \cdot N_{n-1}$. By unfolding this recurrence, we see that, for each $n \ge 0$,

$$N_n = 3 \cdot 4^n.$$

• For any integer $n \geq 0$, let ℓ_n be the length of one single edge of SF_n . Determine the value of ℓ_n , by deriving a recurrence relation and solving it.

Solution: Since each edge of the triangle SF_0 has length 1, we have $\ell_0 = 1$. Let $n \ge 1$. By construction, the length of each edge in SF_n is one-third of the edge-length in SF_{n-1} . This implies that $\ell_n = \frac{1}{3} \cdot \ell_{n-1}$. By unfolding this recurrence, we see that, for each $n \ge 0$,

$$\ell_n = (1/3)^n.$$

• For any integer $n \geq 0$, let L_n be the total length of all edges of SF_n . Prove that

$$L_n = 3 \cdot \left(\frac{4}{3}\right)^n.$$

Solution: For any $n \geq 0$, we have

$$L_n = N_n \cdot \ell_n$$

= $3 \cdot 4^n \cdot (1/3)^n$
= $3 \cdot (4/3)^n$.

• Let a_0 be the area of the triangle SF_0 . For any integer $n \geq 1$, let a_n be the area of one single triangle that is added when constructing SF_n from SF_{n-1} . Determine the value of a_n , by deriving a recurrence relation and solving it.

Solution: According to the construction, each triangle that is added is equilateral. In high school, you have learned that the height of an equilateral triangle with sides of length ℓ is equal to

$$\frac{1}{2}\ell\sqrt{3}$$
.

(In case you forgot, you can either use Pythagoras to prove this, or you use the fact that the height of such a triangle is equal to $\ell \cdot \sin(\pi/3)$.) Thus, the area of this triangle is equal to

$$\frac{1}{2} \cdot \ell \cdot \frac{1}{2} \ell \sqrt{3},$$

which is a constant times ℓ^2 .

Let $n \geq 1$. Since $\ell_n = \frac{1}{3} \cdot \ell_{n-1}$, it follows that $a_n = \frac{1}{9} \cdot a_{n-1}$. By unfolding this recurrence, we see that, for each $n \geq 0$,

$$a_n = (1/9)^n \cdot a_0.$$

• For any integer $n \geq 1$, let A_n be the total area of all triangles that are added when constructing SF_n from SF_{n-1} . Prove that

$$A_n = \frac{3}{4} \cdot \left(\frac{4}{9}\right)^n \cdot a_0.$$

Solution: Let $n \geq 1$. When constructing SF_n , we add triangles to SF_{n-1} ; each such triangle has area a_n . How many such triangles do we add: We add one triangle for each edge of SF_{n-1} . Since SF_{n-1} has N_{n-1} edges, we get

$$A_n = N_{n-1} \cdot a_n$$

$$= (3 \cdot 4^{n-1}) \cdot \left(\left(\frac{1}{9} \right)^n \cdot a_0 \right)$$

$$= \frac{3}{4} \cdot \left(\frac{4}{9} \right)^n \cdot a_0.$$

• Let $n \geq 1$ be an integer. Prove that the total area of SF_n is equal to

$$\frac{a_0}{5} \cdot \left(8 - 3 \cdot \left(\frac{4}{9}\right)^n\right).$$

Hint: For any real number $x \neq 1$,

$$\sum_{k=1}^{n} x^{k} = x \cdot \frac{1 - x^{n}}{1 - x}.$$

Solution: For $n \geq 0$, let $area_n$ denote the total area of SF_n . Then $area_0 = a_0$ and, for $n \geq 1$,

$$area_n = area_{n-1} + A_n.$$

Note that we have determined A_n above. The question asks to prove that

$$area_n = \frac{a_0}{5} \cdot \left(8 - 3 \cdot \left(\frac{4}{9}\right)^n\right).$$

One way to prove this it to use induction; the recurrence is used in the induction step; in this way, you do not need the hint. Another way is to unfold the recurrence. If you do this, you will get, for each $n \ge 0$,

$$area_n = area_0 + \sum_{k=1}^n A_k = a_0 + \sum_{k=1}^n A_k.$$

This gives us

$$area_n = a_0 + \sum_{k=1}^n \frac{3}{4} \cdot \left(\frac{4}{9}\right)^k \cdot a_0$$
$$= a_0 + \frac{3}{4} \cdot a_0 \sum_{k=1}^n \left(\frac{4}{9}\right)^k.$$

Using the hint with x = 4/9, we get

$$\sum_{k=1}^{n} \left(\frac{4}{9}\right)^{k} = \frac{4}{9} \cdot \frac{1 - (4/9)^{n}}{1 - 4/9}$$
$$= \frac{4}{5} \left(1 - (4/9)^{n}\right).$$

Thus,

$$area_n = a_0 + \frac{3}{4} \cdot a_0 \cdot \frac{4}{5} \left(1 - (4/9)^n \right).$$

After some algebra, you will see that the right-hand side is exactly the value that we are trying to obtain.

Remark: Consider the "limit snowflake", when n goes to infinity. Imagine this snowflake to be a country. Since

$$\lim_{n\to\infty} L_n = \infty,$$

the length of this country's border is infinite. But, since

$$\lim_{n\to\infty} area_n = 8a_0/5,$$

this country has a finite area.