

Carleton University  
COMP 3803, Fall 2025, Test 3

Thursday November 20, 2025

STUDENT NAME:

STUDENT NUMBER:

**40 marks total**

**Question 1:** (10 marks) Consider the pushdown automaton with tape alphabet  $\Sigma = \{a, b\}$ , stack alphabet  $\Gamma = \{\$, S\}$ , set of states  $Q = \{q_0, q_1, q_2\}$ , start state  $q_0$ , and instructions

$$\begin{array}{lll} q_0 a \$ \rightarrow q_1 R \$ S & q_1 a S \rightarrow q_2 R \epsilon & q_2 a \$ \rightarrow q_2 R \$ \\ q_0 b \$ \rightarrow q_0 R \$ & q_1 b S \rightarrow q_0 R \epsilon & q_2 b \$ \rightarrow q_2 R \$ \\ q_0 \square \$ \rightarrow q_0 N \$ & q_1 \square S \rightarrow q_1 N S & q_2 \square \$ \rightarrow q_2 N \epsilon \end{array}$$

Recall that, at the start of the computation, the stack contains the symbol  $\$$  (and nothing else).

Exactly one of the following four strings is accepted by this pushdown automaton.

- $\epsilon$
- $baba$
- $babaab$
- $abba$

Which one is it? Just give the answer. No explanation is needed.

**Question 2:** (10 marks) Let  $L$  be the language

$$L = \{(ba)^n a^n : n \geq 0\}.$$

Give an **informal** description of a pushdown automaton that accepts the language  $L$ . (Describe the algorithm in plain English. Do not give the instructions.)



**Question 3:** (10 marks) Let  $L$  be the language

$$L = \{ww : w \in \{a, b\}^*\}.$$

We want to apply the Pumping Lemma for context-free languages to prove that  $L$  is not context-free: As always, we assume that  $L$  is context-free. Let  $p$  be the pumping length, as given by the Pumping Lemma.

Can we use the string  $s = a^pba^pb$  to obtain a contradiction?

Justify your answer.

**Question 4:** (10 marks) Let  $L$  be the language

$$L = \{a^m b^n c^k : m \geq 0, n \geq 0, k = \max(m, n)\}.$$

Prove that  $L$  is not a context-free language.

EXTRA PAGE FOR ANSWERS