

GHOST CHIMNEYS*

DAVID CHARLTON^{†,¶}, ERIK D. DEMAINE^{||} and MARTIN L. DEMAINE^{**}
Computer Science and Artificial Intelligence Laboratory, MIT, Cambridge, USA
[¶]dcharlton@gmail.com
^{||}edemaine@mit.edu
^{**}mdemaine@mit.edu

VIDA DUJMOVIĆ^{‡,††} and PAT MORIN^{‡,‡‡}
School of Computer Science, Carleton University, Ottawa, Canada
^{††}vida@scs.carleton.ca
^{‡‡}morin@scs.carleton.ca

RYUHEI UEHARA[§]
School of Information Science, JAIST, Ishikawa, Japan
uehara@jaist.ac.jp

Received 19 April 2011
Revised 26 March 2012
Communicated by Gill Barequet

ABSTRACT

A planar point set S is an (i, t) set of ghost chimneys if there exist lines H_0, H_1, \dots, H_{t-1} such that the orthogonal projection of S onto H_j consists of exactly $i + j$ distinct points. We give upper and lower bounds on the maximum value of t in an (i, t) set of ghost chimneys, showing that it is linear in i .

Keywords: Crossing lemma; orthogonal projection; shadow sculpture.

1. Introduction

Once upon a time in Japan, there was a power plant with four chimneys called “ghost chimneys” (obake entotsu, or お化け煙突); see Fig. 1. Although these chimneys were dismantled in 1964, they are still famous in Japan, with toys, books,

*Preliminary version was presented at the 22nd Canadian Conference on Computational Geometry.

[†]Supported in part by Akamai Presidential Fellowship.

[‡]Research supported in part by NSERC.

[§]Supported in part by Ogasawara Foundation for the Promotion of Science and Engineering.



Fig. 1. The ghost chimneys in Japan, part of the Senju Thermal Power Station (1926–1964) maintained by the Tokyo Electric Power Company. [Used with permission from Adachi City.]



Fig. 2. Hello Kitty Obake-Entotsu version (© 1976, 2011 SANRIO CO., LTD. APPROVAL NO. S520196)

manga, and movies referencing them (Fig. 2). They are considered a kind of symbol of industrialized Japan in the old, good age of the Showa era.¹

One of the reasons why they are famous and are called “ghost chimneys” is that they could be seen as two chimneys, three chimneys, or four chimneys depending on the point of view. This phenomenon itself was an accident, but it raises several natural questions. What interval of integers can be realized by such chimneys? How many chimneys do we need to realize the interval? How can we arrange the chimneys to realize the interval?

More precisely, we consider the following problem: given an integer i , what is the maximum value $t(i)$ such that there exists a set of points $S \subset \mathbb{R}^2$ and a set $H_0, H_1, \dots, H_{t(i)-1}$ of lines where, for each $j \in \{0, 1, \dots, t(i) - 1\}$, the orthogonal projection of S onto H_j consists of exactly $i + j$ distinct points? We prove the following result:

Theorem 1. *For any integer $i \geq 1$, $2i \leq t(i) < 123.33i + 1$.*

In addition to Theorem 1, we show that $t(1) = 2$, $t(2) = 5$, $t(3) = 9$, and $12 \leq t(4) \leq 15$. These results show that neither the lower bound nor the upper bound of Theorem 1 is tight for all values of i . Theorem 1 is an immediate consequence of Lemma 1 and Lemma 4, which we prove in the next two sections, respectively.

These ghost-chimney problems relate more generally to understanding what orthogonal projections a single 2D or 3D shape can have. In 2D, some closely related problems have been considered.^{2,3} Past explorations into structures in 3D, known variously as 3D ambigrams, trip-lets, and shadow sculptures, have focused on precise, usually connected projections.^{4,5} Our work was originally motivated

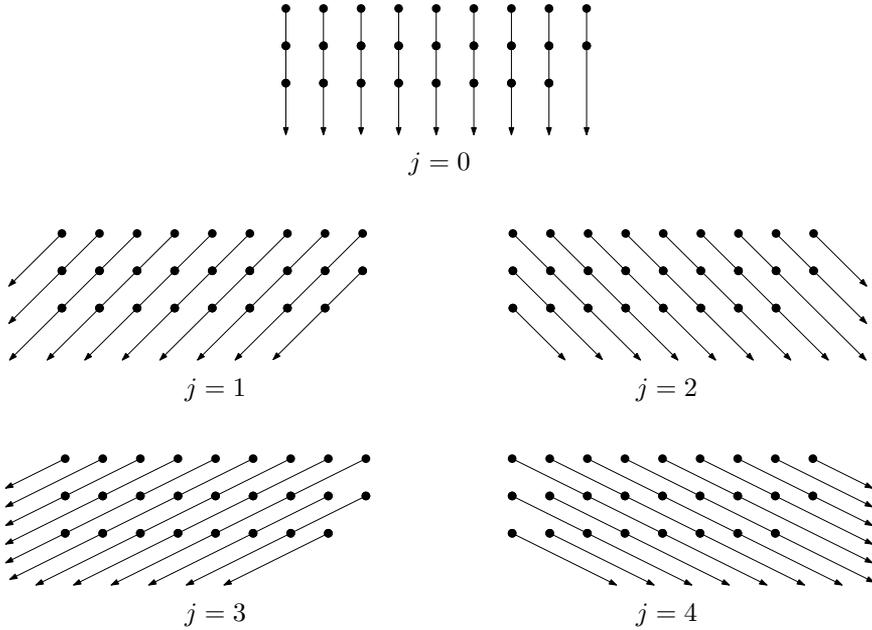


Fig. 3. The set $S(i)$ for $i = 9$ and the projection directions that yield $i, i + 1, \dots, i + 4$ distinct points.

by considering what happens with disconnected projections of unspecified relative position.

2. The Lower Bound

Lemma 1. *For each integer $i \geq 1$, there exists a set $S = S(i)$ of $3i - 1$ points and a set $H_0, H_1, \dots, H_{2i-1}$ of lines such that, for each $j \in \{0, 1, \dots, 2i - 1\}$, the orthogonal projection of S onto H_j has exactly $i + j$ distinct values.*

Proof. The point set S consists of the points of an $i \times 3$ grid with the bottom-right corner removed; see Fig. 3. For even j , H_j is a line of slope $j/2$. For odd j , H_j is a line of slope $-(j + 1)/2$. □

3. The Upper Bound

Our upper-bound proof is closely related to Székely’s proof of the Szemerédi-Trotter Theorem.⁶ We make use of the following version of the Crossing Lemma, which was proved by Pach, Radoičić, Tardos, and Tóth:⁷

Lemma 2 (Crossing Lemma). *Let $\beta = 103/6$, $\gamma = 1024/31827$, and let G be a graph with no self loops, no parallel edges, v vertices, and $e > \beta v$ edges. Then the*

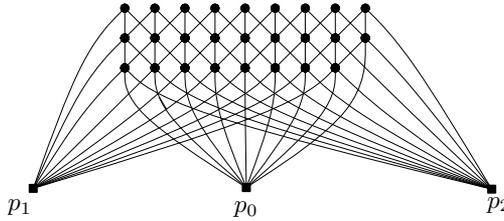


Fig. 4. The graph G for a set of points with $i = 9$ and $t = 3$.

crossing number $\text{cr}(G)$, the minimum number of edge crossings in a certain drawing of G , is given by

$$\text{cr}(G) \geq \gamma \cdot \frac{e^3}{v^2}.$$

Lemma 3. *Let S be a set of r points, and let H_0, H_1, \dots, H_{t-1} be a set of lines such that the orthogonal projection of S onto H_j gives exactly $i + j$ distinct values. Then, $t \leq 34$ or $r \leq (2i/t + 2 + t/2i)i/\gamma$.*

Proof. Each projection direction H_j defines a set L_j of $i + j$ parallel lines, each of which contains at least one point of S . Let G be the geometric graph that contains the points in S plus t additional points p_0, p_1, \dots, p_{t-1} . Two vertices in S are connected by an edge in G if and only if they occur consecutively on some line in $L = \bigcup_{j=0}^{t-1} L_j$. Additionally, each vertex p_j is connected to each of the $i + j$ lexically largest points on each of the lines in L_j . See Fig. 4.

The graph G has $t + r$ vertices and tr edges. Observe that we have a drawing of G so that the only crossings between edges occur where lines in L intersect each other. The total number X of intersecting pairs of lines in L is

$$\begin{aligned} X &\leq \sum_{j=1}^{t-1} (i + j) \cdot \sum_{k=0}^{j-1} (i + k) \leq \sum_{j=1}^{t-1} (i + j)(ij + j^2/2) \\ &\leq \sum_{j=1}^{t-1} (i^2j + 3ij^2/2 + j^3/2) \leq i^2t^2/2 + it^3/2 + t^4/8. \end{aligned}$$

Applying Lemma 2, we learn that either

$$tr \leq \beta(t + r), \tag{1}$$

or

$$X \geq \text{cr}(G) \geq \gamma \frac{(tr)^3}{(t + r)^2}. \tag{2}$$

In the former case, we rewrite (1) to obtain

$$t \leq \beta(t/r + 1) \leq 2\beta \leq 34 + 1/3,$$

hence $t \leq 34$ (since t is an integer).

In the latter case, we expand (2) to obtain

$$i^2t^2/2 + it^3/2 + t^4/8 \geq \gamma \frac{(tr)^3}{(t+r)^2}.$$

Thus

$$i \left(\frac{i}{2t} + \frac{1}{2} + \frac{t}{8i} \right) \geq \gamma \frac{r^3}{(t+r)^2} \geq \gamma r/4,$$

where the second inequality follows from the fact that $t \leq i + t - 1 \leq r$. Rewriting to isolate r finally gives

$$r \leq \left(\frac{2i}{t} + 2 + \frac{t}{2i} \right) i/\gamma,$$

which completes the proof. □

Lemma 4. *For all integers $i \geq 1$, $t(i) < 123.33i + 1$.*

Proof. The existence of H_0 and H_1 implies that the points of S lie on the intersection of i parallel lines with another set of $i + 1$ parallel lines. Thus, $|S| \leq i(i + 1)$, so $t(i) \leq |S| - i + 1 \leq i^2 + 1$. Using this inequality $t(i) \leq i^2 + 1$, the lemma follows for $1 \leq i \leq 123$.

We assume that $i > 123$. By Lemma 3, we have $t \leq 34$ or $r \leq (2i/t + 2 + t/2i)i/\gamma$. When $t \leq 34$, we are done, so suppose $t > 34$. Observe that, if the conditions of Lemma 3 hold for t , then they also hold for all values $t' \in \{35, \dots, t\}$. Therefore the statement $r \leq (2i/t' + 2 + t'/2i)i/\gamma$ is true for all $35 \leq t' \leq t$. If $t \leq 2i$, then we are done. Otherwise, taking $t' = 2i$ yields

$$r \leq (2i/t' + 2 + t'/2i)i/\gamma \leq 4i/\gamma .$$

Combining this with $i + t - 1 \leq r$, we have $i + t - 1 \leq 4i/\gamma$, or $t < 123.33i + 1$. □

4. Small Values of i

In this section we give some tighter bounds on $t(i)$ for $i \in \{1, 2, 3, 4\}$.

Lemma 5. $t(1) = 2$, and $t(2) = 5$.

Proof. Point sets achieving these bounds are the 1×2 and the 2×3 grid, respectively; see Fig. 5. That these point sets are optimal follows from the inequality $t(i) \leq i^2 + 1$ shown in the proof of Lemma 4. □

Notice that the proof of Lemma 5 implies that, for any i , $t(i) \leq i^2 + 1$. The following lemma shows that, for $i \geq 3$, $t(i) \leq i^2$. Of course, this upper bound is tighter than Lemma 4 for $i \leq 123$.

Lemma 6. $t(3) = 9$.

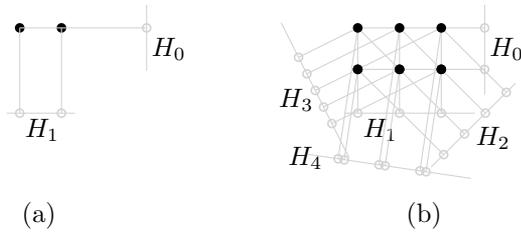


Fig. 5. Point sets showing that (a) $t(1) \geq 2$ and (b) $t(2) \geq 5$.

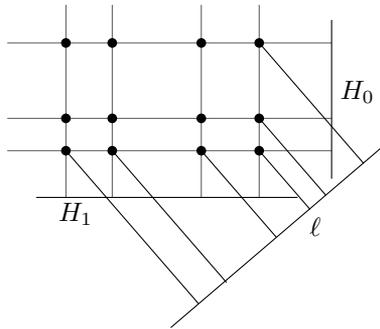


Fig. 6. The proof of Lemma 6.

Proof. The point set $S(4)$ described in the proof of Lemma 1 results in 3 distinct points when they are projected onto a vertical line, therefore $t(3) \geq 9$.

For the upper bound, refer to Fig. 6. By an affine transformation, we may assume that H_0 is vertical and H_1 is horizontal. Thus, the points of S are contained in the intersection of 3 horizontal lines with 4 vertical lines. This establishes that $|S| \leq 12$, so $t(3) \leq 10$. To see that $|S| < 12$, assume otherwise and consider any line ℓ that is neither horizontal nor vertical. By a reflection through a horizontal line, we may assume that ℓ has positive slope, so that every point on the bottom row and right column of S has a distinct projection onto ℓ , so S projects onto at least 6 distinct points on ℓ . In particular, this implies that there is no line H_2 such that S projects onto 5 distinct points on H_2 . \square

Lemma 7. $12 \leq t(4) \leq 15$.

Proof. The point set and lines H_0, H_1, \dots, H_{10} that show $t(4) \geq 12$ are shown in Fig. 7. (H_{11} is omitted since any sufficiently general line will do.)

To see that $t(4) \leq 15$, we argue as in the proof of the second half of Lemma 6. This establishes that $|S| \leq 20$. If $|S| \in \{19, 20\}$ then the number of distinct projections of S onto ℓ is at least 7, but this contradicts the existence of H_2 . Thus, we must have $|S| \leq 18$, and hence $t(4) \leq 15$. \square

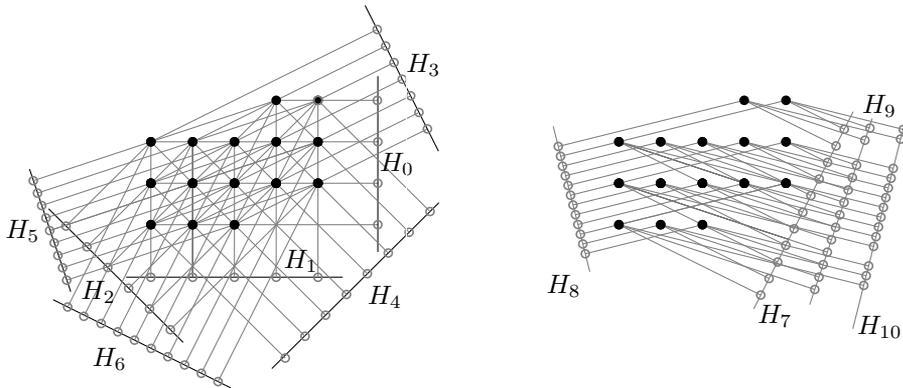


Fig. 7. A $(4, 12)$ set of ghost chimneys.

5. Conclusions

We have given upper and lower bounds on the largest possible value of t , as a function of i , in an (i, t) set of ghost chimneys. These bounds differ by only an (admittedly large) constant factor. Reducing this factor remains an open problem. For small values of i , we have shown that $t(1) = 2$, $t(2) = 5$, $t(3) = 9$, and $12 \leq t(4) \leq 15$.

Another open problem is the generalization of these results to three, or more, dimensions. Given an integer i , what is the maximum value $t(i)$ such that there exists a set of points $S \subset \mathbb{R}^d$ and a set $H_0, H_1, \dots, H_{t(i)-1}$ of hyperplanes where, for each $j \in \{0, 1, \dots, t(i) - 1\}$, the orthogonal projection of S onto H_j consists of exactly $i + j$ distinct points?

Acknowledgments

This work was initiated at the 25th Bellairs Winter Workshop on Computational Geometry, co-organized by Erik Demaine and Godfried Toussaint, held on February 6–12, 2010, in Holetown, Barbados. We thank the other participants of that workshop for providing a stimulating research environment. The authors also thank to Adachi city for their permission to use the photos including one in Fig. 1.

References

1. Adachi City, お化け煙突 (Ghost chimneys, in Japanese), www.city.adachi.tokyo.jp/003/d10100040.html (2009).
2. S. Skiena, Counting k -projections of a point set, *J. Combinat. Theor. Ser. A* **55**(1) (1990) 153–160.
3. J. Matoušek, A. Přívětivý and P. Škovroň, How many points can be reconstructed from k projections? *SIAM J. Discr. Math.* **22**(4) (2008) 1605–1623.
4. J. Keiren, F. van Walderveen and A. Wolff, Constructability of trip-lets, *Abstracts from the 25th European Workshop on Comp. Geom.*, Brussels, Belgium (2009).

5. N. J. Mitra and M. Pauly, Shadow art, *ACM Trans. Graphics* **28**(5) (2009) Article 157.
6. L. A. Székely, Crossing numbers and hard Erdős problems in discrete geometry, *Combinat. Prob. Comput.* **6** (1997) 353–358.
7. J. Pach, R. Radoičić, G. Tardos and G. Tóth, Improving the crossing lemma by finding more crossings in sparse graphs, *Proc. 20th Annual Symp. Comp. Geom.* (2004), pp. 68–75.